

Max Heap: Representation

- Since the heap is a complete binary tree, we could adopt "Array Representation" as we mentioned before!
- Let node i be in position i (array[0] is empty)
 - $Parent(i) = \lfloor i/2 \rfloor$ if $i \neq 1$. If i = 1, i is the root and has no parent.
 - leftChild(i) = 2i if $2i \le n$. If 2i > n, then i has no left child.
 - rightChild(i) = 2i + 1 if $2i + 1 \le n$, if 2i + 1 > n, then i has no right child.

Max Heap: Insert

- Make sure it is a complete binary tree
- Insert a new node
- Check if the new node is greater than its parent
- If so, swap two nodes



Max Heap: Delete

- I. Always delete the root
- 2. Move the last element to the root (maintain a complete binary tree)
- 3. Swap with larger and largest child (if any)
- 4. Continue step 3 until the max heap is maintained (trickle down)



Place Proposed to the max-heap structure

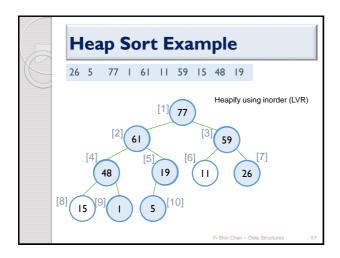
Utilize the max-heap structure

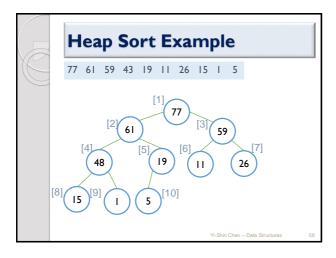
The insertion and deletion could be done in O(logn)

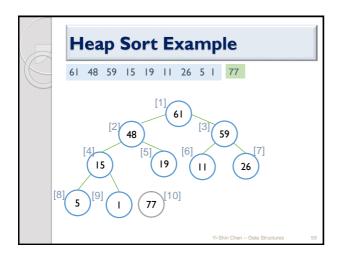
Build a max-heap using n records, insert each record one by one (O(nlogn))

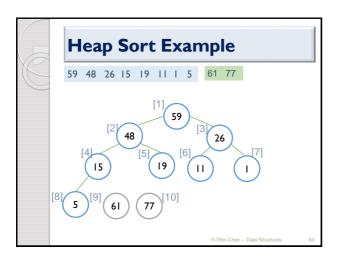
Iteratively remove the largest record (the root) from the max-heap (O(nlogn))

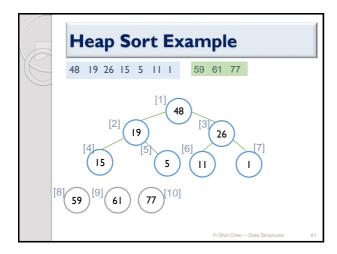
Not a stable sort

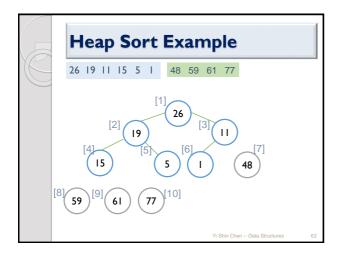


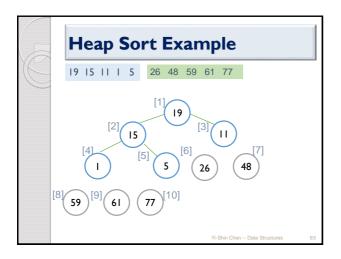


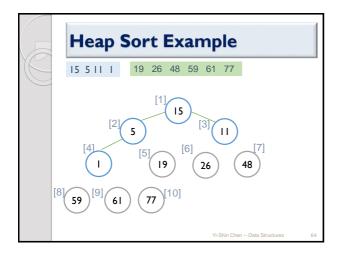


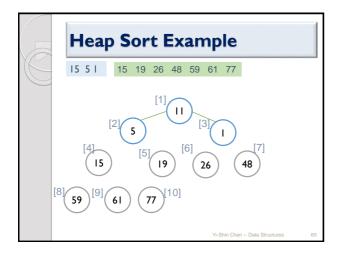


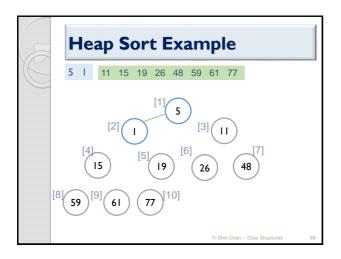












	Heap Sort Example
	1 5 11 15 19 26 48 59 61 77
	[2] (5) [1] [3] [1]
	$ \begin{bmatrix} 4 \\ \hline 15 \end{bmatrix} \begin{bmatrix} 5 \\ \hline 19 \end{bmatrix} \begin{bmatrix} 6 \\ 26 \end{bmatrix} \begin{bmatrix} 7 \\ 48 \end{bmatrix} \begin{bmatrix} 7 \\ 48 \end{bmatrix} $
[8]	[9] 6I 77 [10]
	Yi-Shin Chen Data Structures 67